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## Principles of Distributed Computing Exercise 4

## 1 Bad Queues in a Mesh

In Section 4.2, Remark 2 states that the greedy algorithm on a mesh might need queues in the order of m (on an  $m \times m$  mesh). Can you find an instance where the queue size gets really big (in the sense of as close to m as possible)?

## 2 Good Queues in a Mesh

Prove that the time complexity for random destination routing is in O(m), see Remark 1 of Theorem 4.10. To that end, show that the probability that a single column contains 2em destination packets is "really small" and finish the proof.